

Admin

- ◇ Assign 2 due Wed
- ◇ Today's topics
 - Functional recursion
- ◇ Reading
 - Reader ch. 4-5-6 (today-W-F)

Lecture #8

Solving problems recursively

- ◇ A recursive function calls itself — wacky!
- ◇ Idea: solve problem using coworkers (clones) who work and act like you
 - Delegate similar, smaller problem to clone
 - Combine result from clone(s) to solve total problem
 - Work toward trivial version that is directly solvable
- ◇ For problems that exhibit "self-similarity"
 - Structure repeats within at different levels of scale
 - Solving larger problem means solving smaller problem(s) within
- ◇ Feels mysterious at first
 - "Leap of faith" required
 - With practice, master the art of recursive decomposition
 - Eventually grok the underlying patterns

Functional recursion

- ◇ Function that returns answer/result
 - Outer problem result uses result from smaller, same problem(s)
- ◇ Base case
 - Simplest version of problem
 - Can be directly solved
- ◇ Recursive case
 - Make call(s) to self to get results for smaller, simpler version(s)
 - Recursive calls must advance toward base case
 - Results of recursive calls combined to solve larger version

Power example

- ◇ C++ has no exponentiation op
- ◇ Iterative formulation for Raise function
 - $\text{base}^{\text{exp}} = \text{base} * \text{base} * \dots * \text{base}$ (exp times)

```
int Raise(int base, int exp)
{
    int result = 1;
    for (int i = 0; i < exp; i++)
        result *= base;
    return result;
}
```

Recursive version

◇ Now consider recursive formulation

- $\text{base}^{\text{exp}} = \text{base} * \text{base}^{\text{exp}-1}$

```
int Raise(int base, int exp)
{
    if (exp == 0) return 1; } Base case
    else return base * Raise(base, exp-1); } Recursive case
}
```

More efficient recursion

$\text{base}^{\text{exp}} = \text{base}^{\text{exp}/2} * \text{base}^{\text{exp}/2}$ (* base if exp is odd)

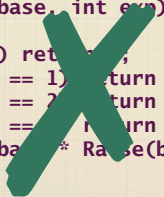
```
int Raise(int base, int exp)
{
    if (exp == 0) return 1;
    else {
        int half = Raise(base, exp/2);
        if (exp % 2 == 0) return half * half;
        else return base * half * half;
    }
}
```

Avoid "arm's length" recursion

◇ Aim for simple, clean base case

- No need to anticipate other earlier stopping points
- Avoid looking ahead before recursive calls, just let simple base case handle

```
int Raise(int base, int exp)
{
    if (exp == 0) return 1;
    else if (exp == 1) return base;
    else if (exp == 2) return base * base;
    else if (exp == 3) return base * base * base;
    else return base * Raise(base, exp - 1);
}
```



Recursion and efficiency

- ◇ Recursion provides no guarantee of (in)efficiency
 - Recursion can require same resources as alternative approach
 - Or recursion may be much more or much less efficient
 - For problems with simple iterative solution, iteration is likely the best
- ◇ Why recursion then?
 - Can express with clear, direct, elegant code
 - Can intuitively model a task that is recursive in nature
 - Solution may require recursion — iteration won't do!

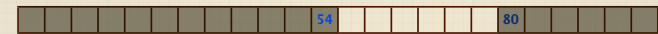
Palindromes

- ◇ A palindrome string reads same when reversed
 - e.g. "was it a car or a cat i saw", "go hang a salami im a lasagna hog"
- ◇ Recursive insight
 - First and last letter match and interior is palindrome
- ◇ Base case?

```
bool IsPalindrome(string s)
{
    if (s.length() <= 1) return true;
    return s[0] == s[s.length()-1] &&
        IsPalindrome(s.substr(1, s.length()-2));
}
```

Binary search

- ◇ Searching for key within vector
 - Linear search starts at beginning and searches to end
 - Binary search uses divide-and-conquer (requires **sorted** vector)
 - Much faster method!
- ◇ Recursive insight:
 - Consider middle elem of vector, if key, you're done
 - Otherwise decide which half to recursively search



- ◇ Base case?

Binary search code

```
const int NotFound = -1;

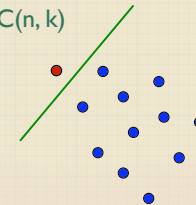
int BSearch(Vector<string> &v,
            int start, int stop, string key)
{
    if (start > stop) return NotFound;

    int mid = (start + stop)/2;
    if (key == v[mid])
        return mid;
    else if (key < v[mid])
        return BSearch(v, start, mid-1, key);
    else
        return BSearch(v, mid+1, stop, key);
}
```

- ◇ Classic "divide and conquer" algorithm
 - Operates very efficiently! Double size of vector, how much longer to search?

Choosing a subset

- ◇ Reader ch 4, exercise 8
 - Given N things, how many different ways can you choose K of them?
 - e.g. given a dorm of 60 people, how many different groups of 4 people can go together to Flicks?
 - N-choose-K, written as $C(n, k)$



Number of subsets that include ● = $C(n-1, k-1)$
+ Number of subsets that don't include ● = $C(n-1, k)$

Choose code

◇ Simplest base case

- when **no** choices remain at all

```
int C(int n, int k)
{
    if (k == 0 || k == n)
        return 1;
    else
        return C(n-1, k) + C(n-1, k-1);
}
```